Vector-partition functions

Matthias Beck

San Francisco State University

math.sfsu.edu/beck

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 – an $(m \times d)$ -integral matrix $\mathbf{b} \in \mathbb{Z}^m$

Goal: Compute vector partition function $\phi_{\mathbf{A}}(\mathbf{b}) := \# \left\{ \mathbf{x} \in \mathbb{Z}_{\geq 0}^d : \ \mathbf{A} \, \mathbf{x} = \mathbf{b} \right\}$

(defined for **b** in the nonnegative linear span of the columns of **A**)

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Applications in...

- ► Number Theory (partitions)
- ▶ Discrete Geometry (polyhedra)
- Commutative Algebra (Hilbert series)
- Algebraic Geometry (toric varieties)
- Representation Theory (tensor product multiplicities)
- Optimization (integer programming)
- ► Chemistry, Biology, Physics, Computer Science, Economics...

An example

$$\mathbf{A} = \left(\begin{array}{cccc} 1 & 2 & 1 & 0 \\ 1 & 1 & 0 & 1 \end{array}\right)$$

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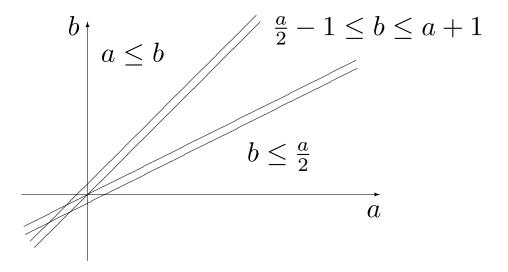
$$\mathbf{A} = \left(\begin{array}{cccc} 1 & 2 & 1 & 0 \\ 1 & 1 & 0 & 1 \end{array}\right)$$

$$\phi_{\mathbf{A}}(a,b) = \# \left\{ \mathbf{x} \in \mathbb{Z}_{\geq 0}^{4} : \mathbf{A} \mathbf{x} = \begin{pmatrix} a \\ b \end{pmatrix} \right\}$$

$$= \begin{cases} \frac{a^{2}}{4} + a + \frac{7 + (-1)^{a}}{8} & \text{if } a \leq b \\ ab - \frac{a^{2}}{4} - \frac{b^{2}}{2} + \frac{a + b}{2} + \frac{7 + (-1)^{a}}{8} & \text{if } \frac{a}{2} - 1 \leq b \leq a + 1 \\ \frac{b^{2}}{2} + \frac{3b}{2} + 1 & \text{if } b \leq \frac{a}{2} \end{cases}$$

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Restricted partition function

$$\phi_{\mathbf{A}}(t) = \# \{ (m_1, \dots, m_d) \in \mathbb{Z}_{>0}^d : m_1 a_1 + \dots + m_d a_d = t \},$$

a quasi-polynomial, i.e., $\phi_{\mathbf{A}}(t)=c_{d-1}(t)\,t^{d-1}+c_{d-2}(t)\,t^{d-2}+\cdots+c_0(t)$ where $c_k(t)$ are periodic

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Frobenius problem: find the largest value for t such that $\phi_{\mathbf{A}}(t) = 0$

Rational (convex) polytope \mathcal{P} – convex hull of finitely many points in \mathbb{Q}^d

Alternative description: $\mathcal{P} = \left\{ \mathbf{x} \in \mathbb{R}^d : \mathbf{A} \mathbf{x} \leq \mathbf{b} \right\}$

For $t \in \mathbb{Z}_{>0}$, let $L_{\mathcal{P}}(t) := \# \left(t \mathcal{P} \cap \mathbb{Z}^d \right)$

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Theorem (Ehrhart 1967) If $\mathcal P$ is a rational polytope, then the functions $L_{\mathcal P}(t)$ and $L_{\mathcal P^\circ}(t)$ are quasi-polynomials in t of degree $\dim \mathcal P$. If $\mathcal P$ has integer vertices, then $L_{\mathcal P}$ and $L_{\mathcal P^\circ}$ are polynomials. Furthermore, $L_{\mathcal P}(0)=1$

Theorem (Ehrhart, Macdonald 1970) $L_{\mathcal{P}}(-t) = (-1)^{\dim \mathcal{P}} L_{\mathcal{P}^{\circ}}(t)$

The computation of the (Ehrhart-)quasi-polynomial

$$\phi_{\mathbf{A}}(t) = \# \{ (m_1, \dots, m_d) \in \mathbb{Z}_{>0}^d : m_1 a_1 + \dots + m_d a_d = t \},$$

gives rise to the Fourier-Dedekind sum (MB-Robins 2003)

$$\sigma_n(a_1, \dots, a_d; a_0) := \frac{1}{a_0} \sum_{\lambda^{a_0} = 1} \frac{\lambda^n}{(1 - \lambda^{a_1}) \cdots (1 - \lambda^{a_d})}.$$

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Choosing $d=2, n=0, a_1=a, a_2=1, a_0=b$ gives rise to the classical Dedekind sum

$$s(a,b) := \frac{1}{4b} \sum_{j=1}^{b-1} \cot\left(\frac{\pi j a}{b}\right) \cot\left(\frac{\pi j}{b}\right)$$

Ehrhart-Macdonald Reciprocity yields the functional identity

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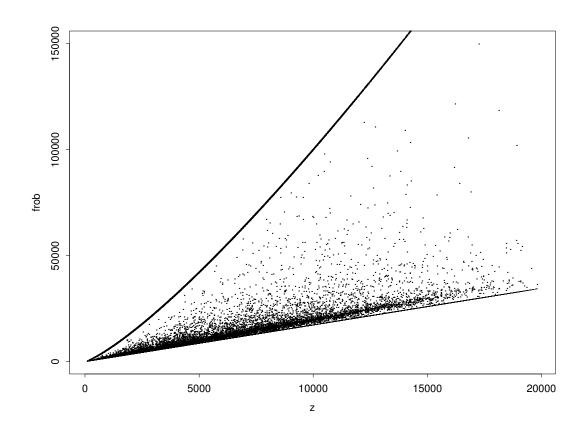
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The identity

$$\phi_{\mathbf{A}}(t) = 0$$
 for $-(a_1 + \dots + a_d) < t < 0$

gives a new reciprocity relation which is a "higher-dimensional" analog of that for the Dedekind-Rademacher sum.

Algorithms, bounds, experimental data on Frobenius problem (MB–Einstein–Zacks 2003)



Shameless plug

M. Beck & S. Robins

Computing the continuous discretely Integer-point enumeration in polyhedra

To appear in Springer Undergraduate Texts in Mathematics

Preprint available at math.sfsu.edu/beck

Vector partition theorems

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Quasi-polynomial – a finite sum $\sum_{\mathbf{n}} c_{\mathbf{n}}(\mathbf{b}) \mathbf{b}^{\mathbf{n}}$ with coefficients $c_{\mathbf{n}}$ that are functions of \mathbf{b} which are periodic in every component of \mathbf{b} .

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A matrix is unimodular if every square submatrix has determinant ± 1 .

Theorem (Sturmfels 1995) $\phi_{\mathbf{A}}(\mathbf{b})$ is a piecewise-defined quasi-polynomial in **b** of degree $d - \operatorname{rank}(\mathbf{A})$. The regions of \mathbb{R}^m in which $\phi_{\mathbf{A}}(\mathbf{b})$ is a single quasi-polynomial are polyhedral. If A is unimodular then ϕ_A is a piecewise-defined polynomial.

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Theorem (MB 2002) Let r_k denote the sum of the entries in the k^{th} row of \mathbf{A} , and let $\mathbf{r} = (r_1, \dots, r_m)$. Then $\phi_{\mathbf{A}}(\mathbf{b}) = (-1)^{d-\operatorname{rank} \mathbf{A}} \phi_{\mathbf{A}}(-\mathbf{b} - \mathbf{r})$

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- ▶ Given one such region, compute the (quasi-)polynomial $\phi_{\mathbf{A}}(\mathbf{b})$
- ▶ Barvinok: $\sum_{t\geq 0} \phi_{\mathbf{A}}(t\mathbf{b}) z^t$ can be computed in polynomial time

► Columns of A – vectors of a classical root system

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- (Baldoni–MB–Cochet–Vergne 200?) Computational approach using Jeffrey-Kirwan residues and DeConcini-Prochesi's maximal nested sets

Euler's generating function

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ight)$$

 $\phi_{\mathbf{A}}(\mathbf{b})$ equals the coefficient of $\mathbf{z}^{\mathbf{b}}:=z_1^{b_1}\cdots z_m^{b_m}$ of the function

$$\frac{1}{(1-\mathbf{z}^{\mathbf{c}_1})\cdots(1-\mathbf{z}^{\mathbf{c}_d})}$$

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Proof Expand each factor into a geometric series.

$$\phi_{\mathbf{A}}(\mathbf{b}) = \operatorname{const} \frac{1}{(1 - \mathbf{z}^{\mathbf{c}_1}) \cdots (1 - \mathbf{z}^{\mathbf{c}_d}) \mathbf{z}^{\mathbf{b}}}$$

Partial fractions

$$\phi_{\mathbf{A}}(\mathbf{b}) = \operatorname{const} \frac{1}{(1 - \mathbf{z}^{\mathbf{c}_1}) \cdots (1 - \mathbf{z}^{\mathbf{c}_d}) \mathbf{z}^{\mathbf{b}}}$$

Expand into partial fractions in z_1 :

$$\frac{1}{(1-\mathbf{z}^{\mathbf{c}_1})\cdots(1-\mathbf{z}^{\mathbf{c}_d})\mathbf{z}^{\mathbf{b}}} = \frac{1}{z_2^{b_2}\cdots z_m^{b_m}} \left(\sum_{k=1}^d \frac{A_k(\mathbf{z},b_1)}{1-\mathbf{z}^{\mathbf{c}_k}} + \sum_{j=1}^{b_1} \frac{B_j(\mathbf{z})}{z_1^j} \right)$$

Here A_k and B_j are polynomials in z_1 , rational functions in z_2, \ldots, z_m , and exponential in b_1 .

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$$= \operatorname{const} \frac{1}{z_{2}^{b_{2}} \cdots z_{m}^{b_{m}}} \sum_{k=1}^{d} \frac{A_{k}(0, z_{2}, \dots, z_{m}, b_{1})}{1 - (0, z_{2}, \dots, z_{m})^{\mathbf{c}_{k}}}$$

Advantages

easy to implement



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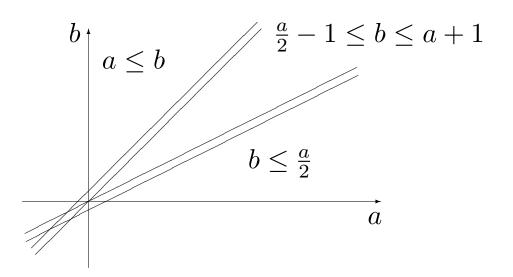
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Advantages

- ► easy to implement
- ▶ allows symbolic computation
- constraints which define the regions of (quasi-)polynomiality are obtained "automatically"

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$$x_1, x_2, x_3, x_4 \ge 0$$

 $x_1 + 2x_2 + x_3 = a$
 $x_1 + x_2 + x_4 = b$

