Proximity Graphs and Principal Curves for Shape Skeletonization

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The Problem:

Investigate the relationship between shape skeletons, proximity graphs, and principal curves.

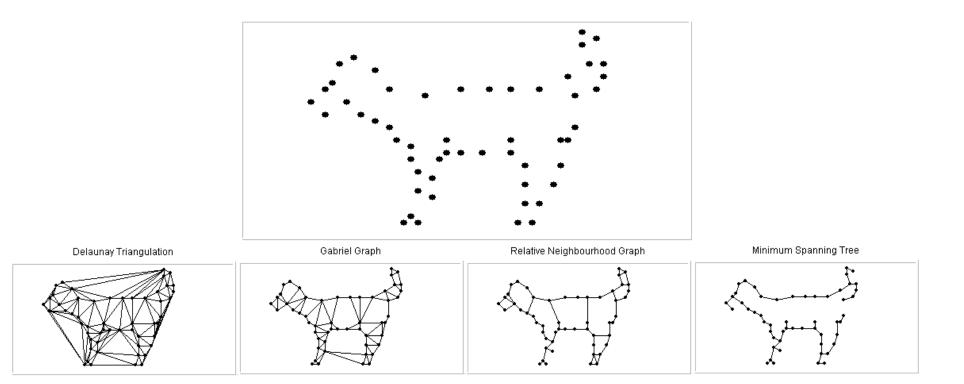
So what is a shape skeleton?

What are Proximity Graphs? (Redux)

- A proximity graph is a simply a graph in which two vertices are connected by an edge if and only if the vertices satisfy particular geometric requirements.
- "Proximity" here means spatial distance.
- Many of these graphs can be formulated with respect to many metrics, but the Euclidean metric is used most frequently.

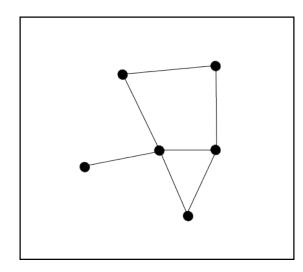
Why Proximity Graphs?

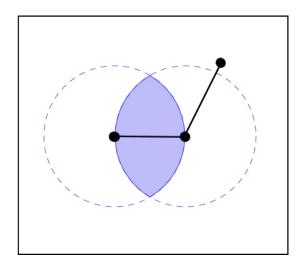
 Proximity graphs have been suggested as a convenient and relatively efficient way of generating a "primal sketch" of the shape of a point set—all of the graphs below can be created in O(n log n) time.



The RNG

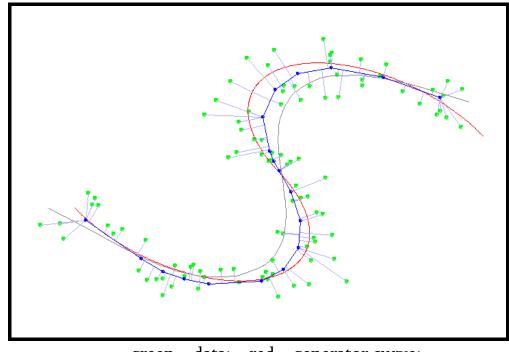
- Let Λ(p,q) be the intersection of the circle about p with a radius of dist(p,q) and the circle about q with a radius of dist(q,p). This is called a *lune*.
- The *relative neighborhood graph*, *RNG*(*V*) of a set of points *V*, is the graph that has an edge (*p*,*q*) if and only if the intersection of Λ(*p*,*q*) and *V* is empty.
- You can think of the *RNG* as a graph which connects each point to its nearest neighbors in "each direction."





Principal Curves

- Principle curves are smooth curves that pass through the "middle" of a set of points (or a distribution)—"continuous curves of a given length which minimize the expected squared distance between the curve and points of the space randomly chosen according to a given distribution." [Kegl, et al., 2000]
- Non-linear generalization of *principal components*.



green = data; red = generator curve; gray = HS principal curve; blue = KKLZ principal curve

Selected Prior Work

- Hastie, Stuetzle, 1989.
 - $O(n^2)$
- Tibshirani, 1992.
- Kegl, Krzyzak, Linder, Zeger, 2000. (k-segments)
 - $O(n^2)$
 - $O(n^{5/3})$ with standard assumptions
 - adding more than one vertex at a time $O(n^{5/3})$ $O(nk \log k)$
 - setting k constant O(nk)
- Singh, Cherkassy, Papanikolopoulos, 2000. (Self-Organizing Map)
- Verbeek, Vlassis, Kröse, 2002 (Local principal components).
 - $O(kn^2)$

My Approach

- Similar to Dr. Singh's approach, I use a proximity graph to initialize the topology of the skeleton. However, I conducted most of my experiments using relative neighborhood graphs instead of minimum spanning trees.
- Similar to the original HS principal curve algorithm, I performed a perpendicular linear regression and projection on the *n* nearest points along the graph to update the position of points on the graph.
- Using an RNG to create the initial primal sketch has two main advantages:
 - It allows the final skeleton to have a much more complicated topological structure than just a curve, including loops, self-intersections, and branches.
 - The RNG of a planar point set can be found in *O*(*n* log *n*) time, so we can get a rough approximation of the shape very quickly.

...and here it is! Or rather, here they are (2 versions).

Problems!

- This algorithm may not converge! In fact, it some cases it definitely doesn't.
- There are several parameters that must be hand-tuned.
 - For version 1 and 2:
 - The number of nearest points along the RNG to consider when doing the local smoothing.
 - The number of iterations must be specified explicitly, since I haven't developed convergence criteria (and convergence may not even happen).
 - The smoothing function must also be specified. (I have deviated slightly from HS here, using a function the gives point near the current point more weight, and points farther away less weight.)
 - In addition, version 2, takes another parameter: the number of steps to be taken along the previous RNG when constructing new "local" RNGs.
- Supposing the algorithm does converge, the resulting graph may need to be trimmed to get rid of edges that are "too long" relative to the rest.

Provisional Complexity

Version 1 (per iteration), for a fixed number m of points along the graph to be used in the local smoothing step.

Initialize the RNG of the point set	$O(n \log n)$
Then, for each iteration:	
1. Find the RNG of the current graph	$O(n \log n)$
2. For each point	$O(nm^2)$
1. Find the <i>m</i> nearest points along the graph. Assuming that the worst	
possible case is a triangular grid of points, and traversing the graph	
using a BFS strategy $O(m^2)$	
(This is very much a back-of-the envelope estimate—I haven't	
worked out the proof in detail; it could be better or worse.)	
2. Regression/projection $O(m)$	
3. Iterate again	

Provisional Complexity

So the algorithm has a *provisional* complexity of $O(n(m^2 + \log n))$. This gives $O(n \log n)$ for fixed *m*, but this is cold comfort, since *m* will need to grow as *n* does in order for the curve to maintain a reasonable level of smoothness.

Also, this does no one any good until some type of convergence can be guaranteed.

References

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